

CE 191: Civil and Environmental Engineering Systems Analysis

LEC 03 : Graphical Solutions to LP

Professor Scott Moura
Civil & Environmental Engineering
University of California, Berkeley

Fall 2014



Graphical Solutions of Linear Programs

Example:

$$\min \quad J = 140x_1 + 160x_2$$

$$\text{s. to} \quad 2x_1 + 4x_2 \leq 28$$

$$5x_1 + 5x_2 \leq 50$$

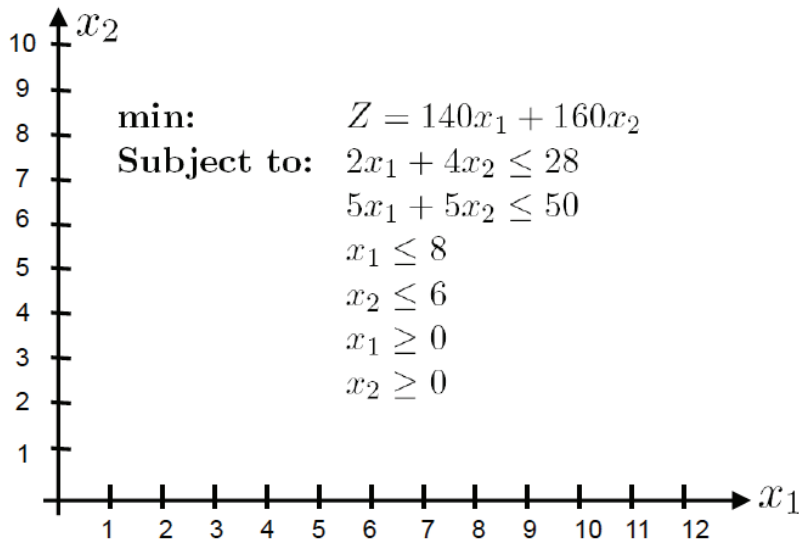
$$x_1 \leq 8$$

$$x_2 \leq 6$$

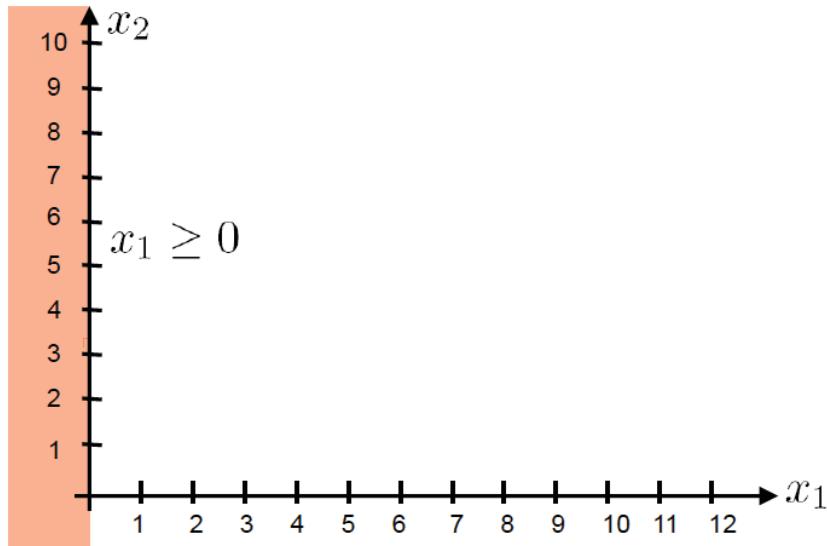
$$x_1 \geq 0$$

$$x_2 \geq 0$$

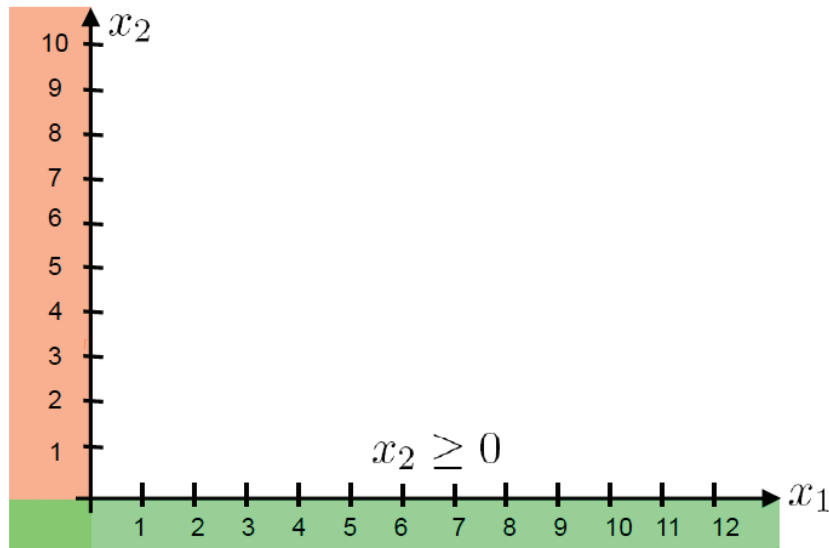
Construction of the feasible set



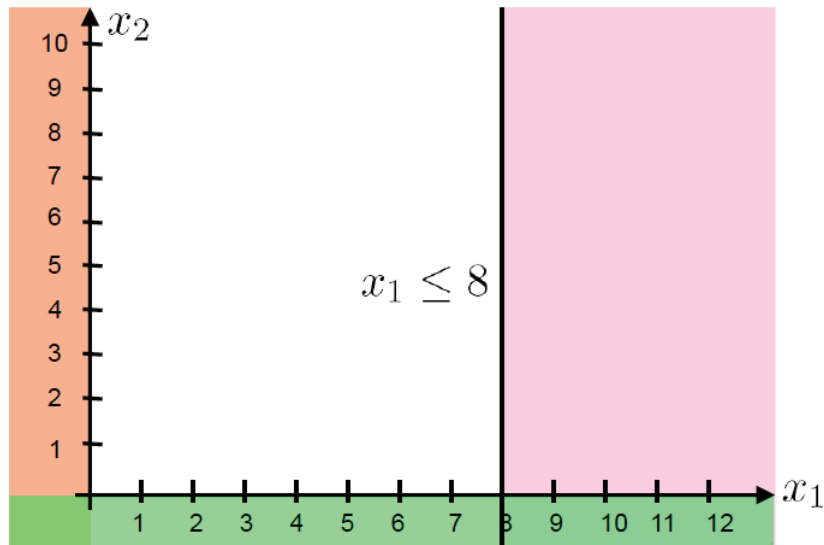
Construction of the feasible set



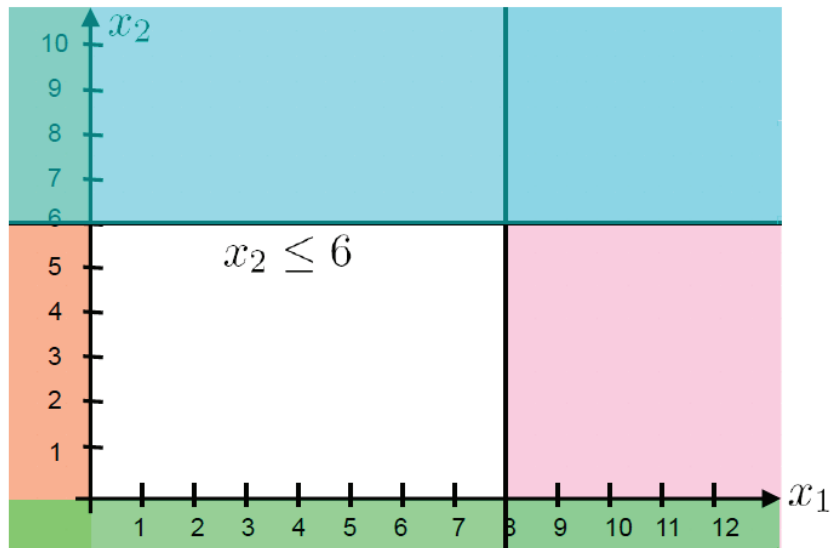
Construction of the feasible set



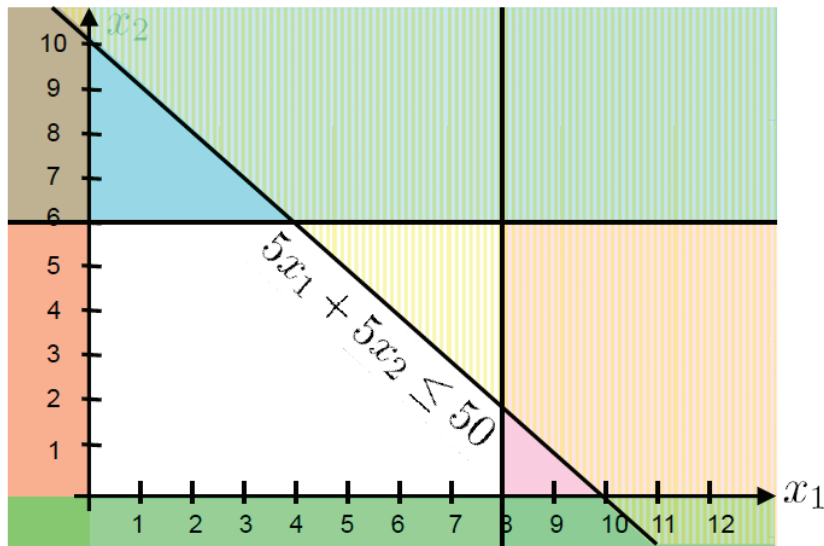
Construction of the feasible set



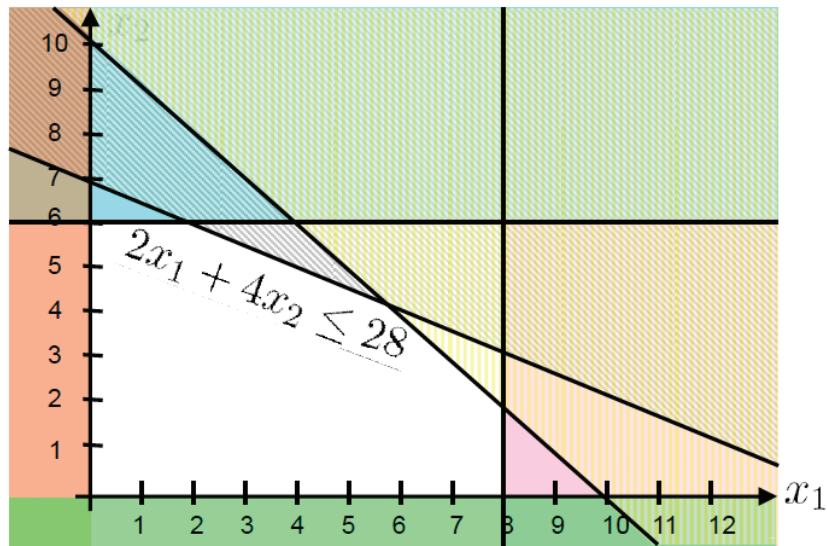
Construction of the feasible set



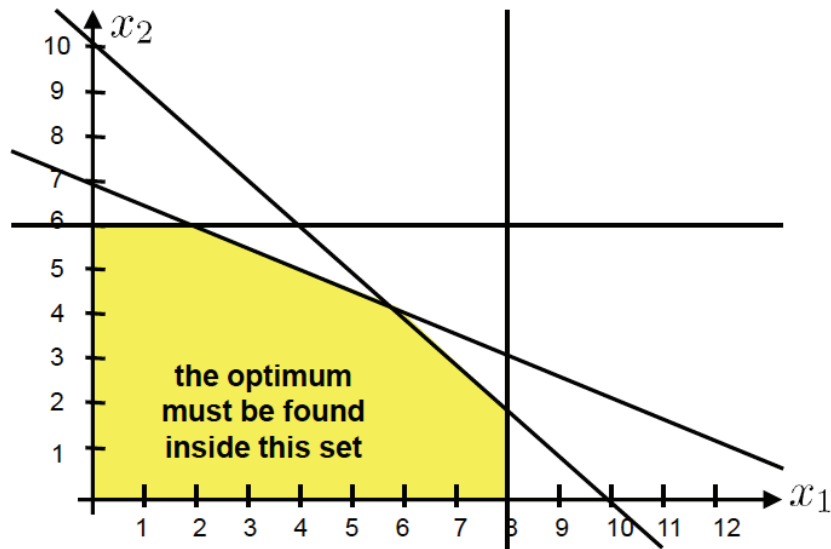
Construction of the feasible set



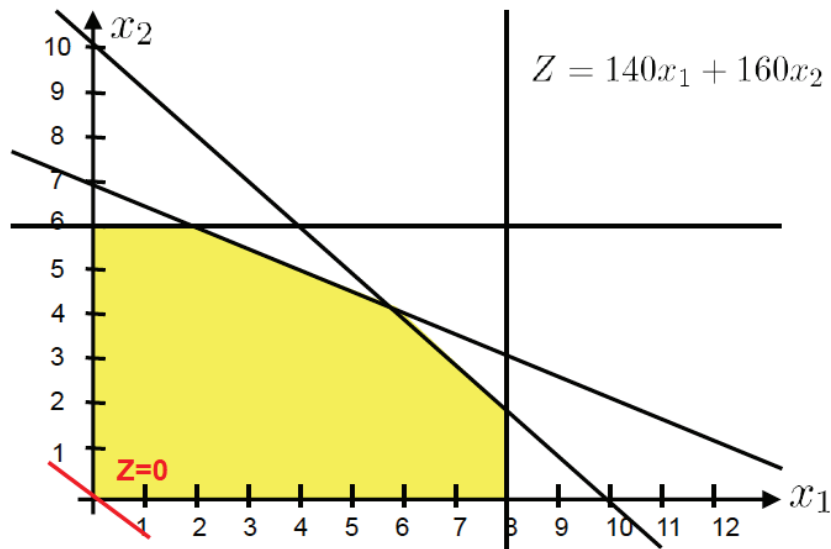
Construction of the feasible set



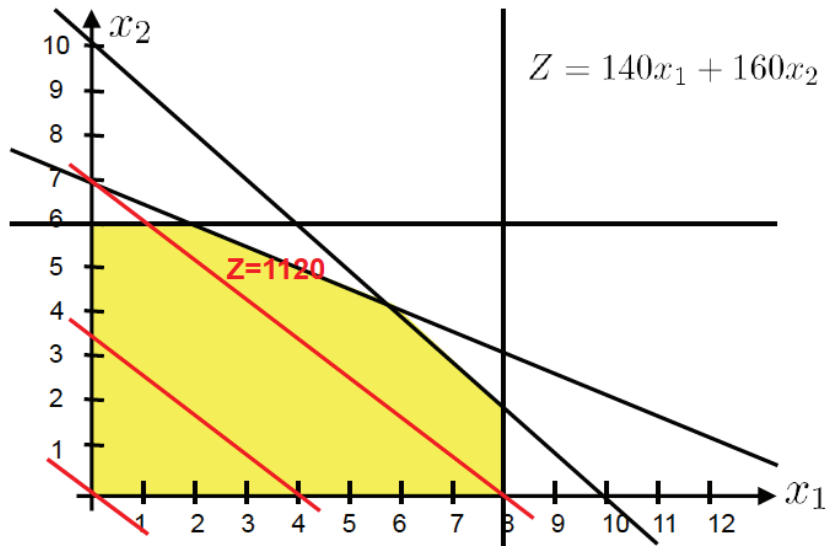
Feasible Set Final Result



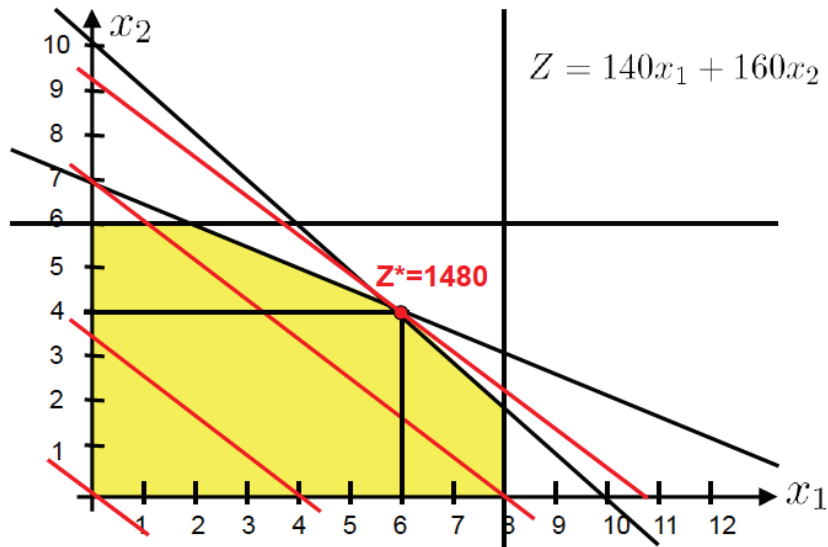
Isolines



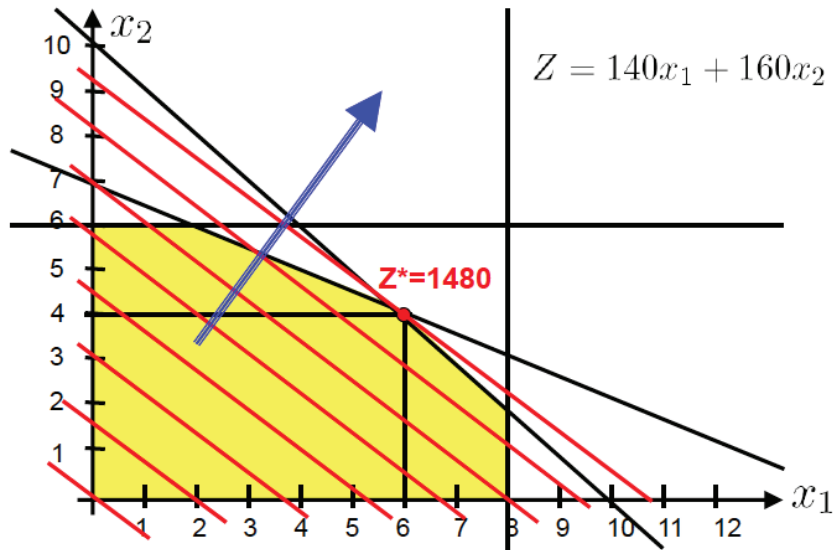
Isolines



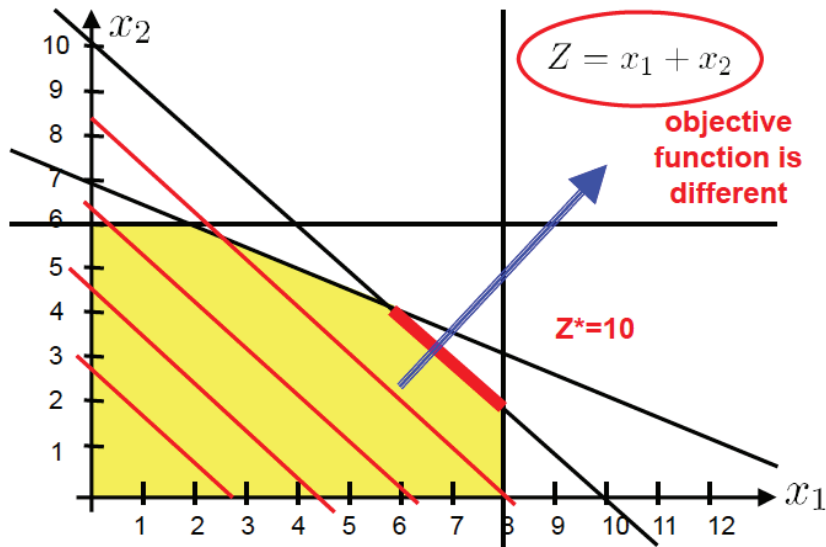
Isolines



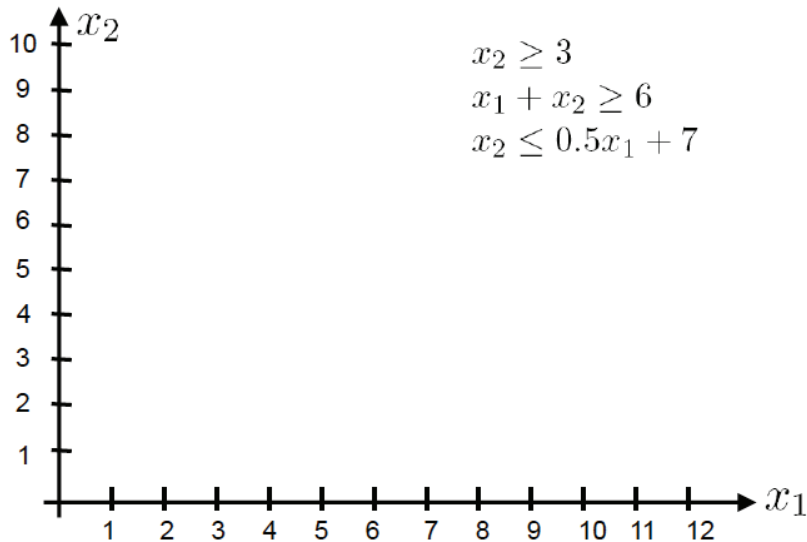
Gradient of the cost function



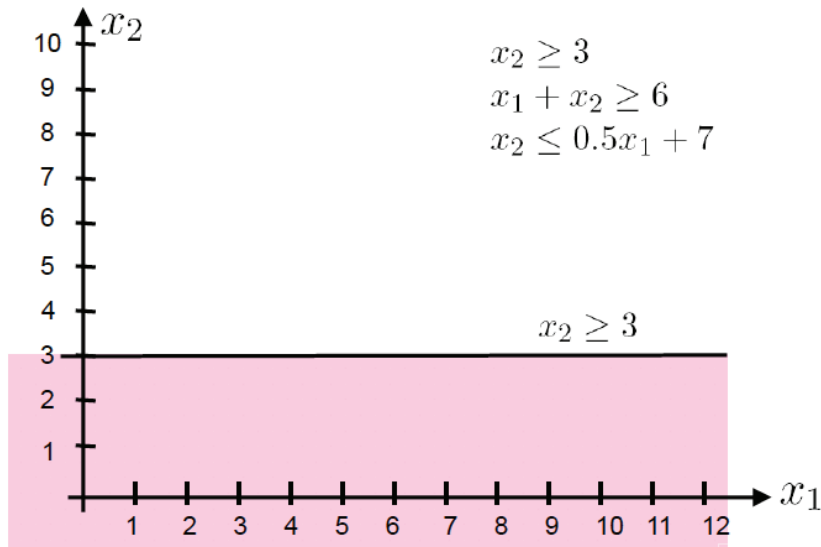
Uniqueness (or not) of the cost function



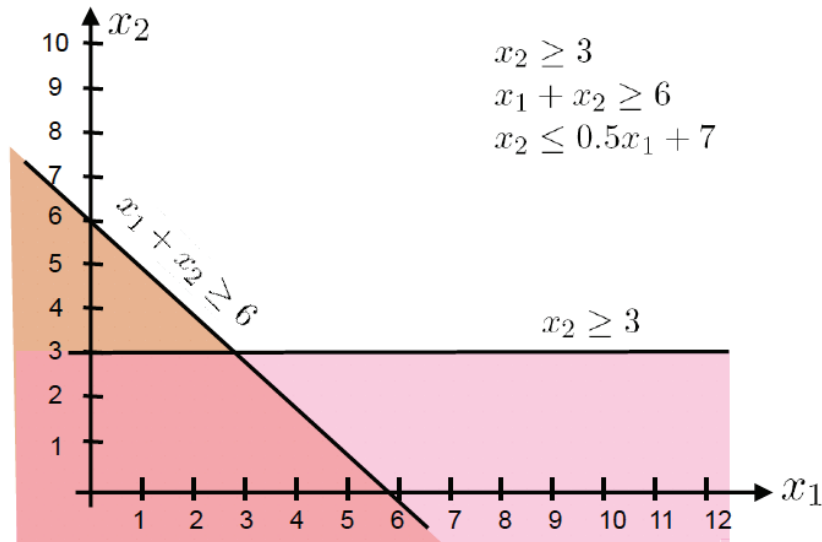
Features of the feasible set



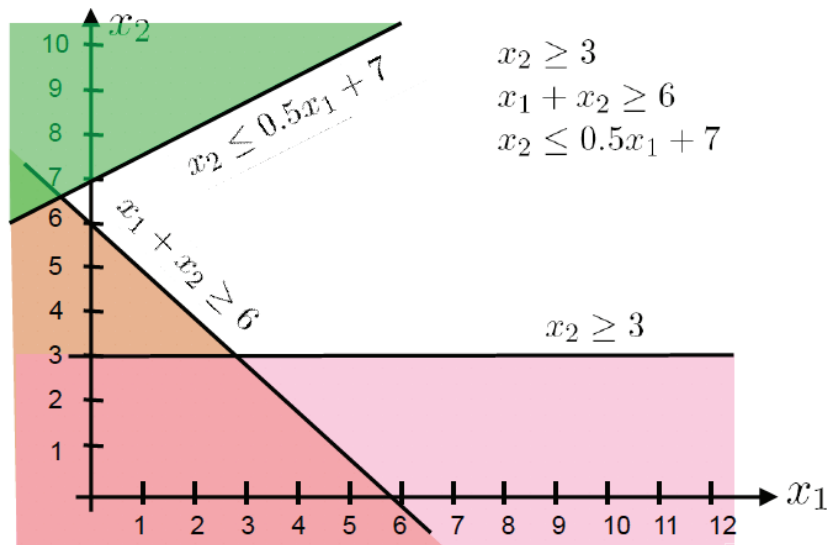
Features of the feasible set



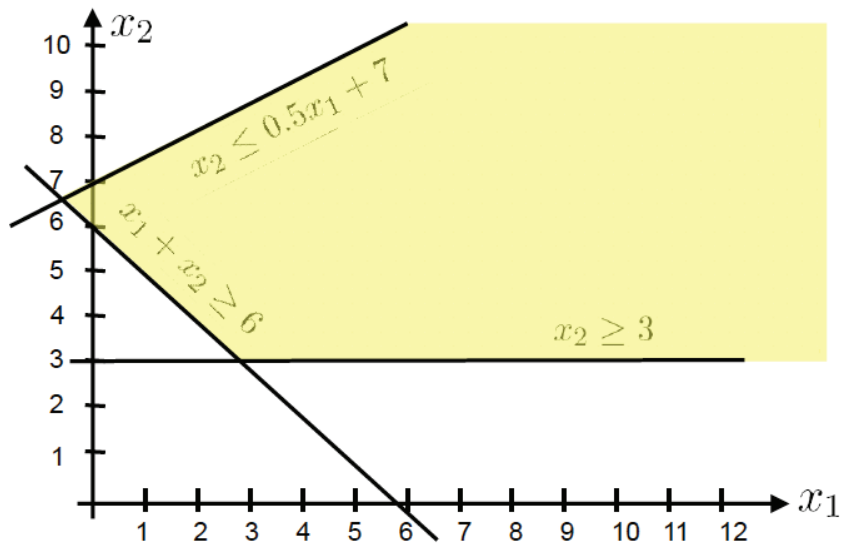
Features of the feasible set



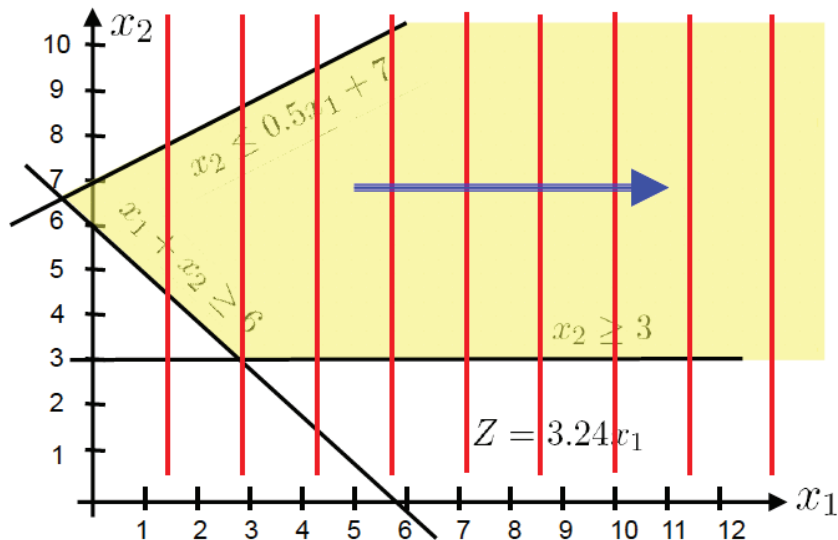
Features of the feasible set



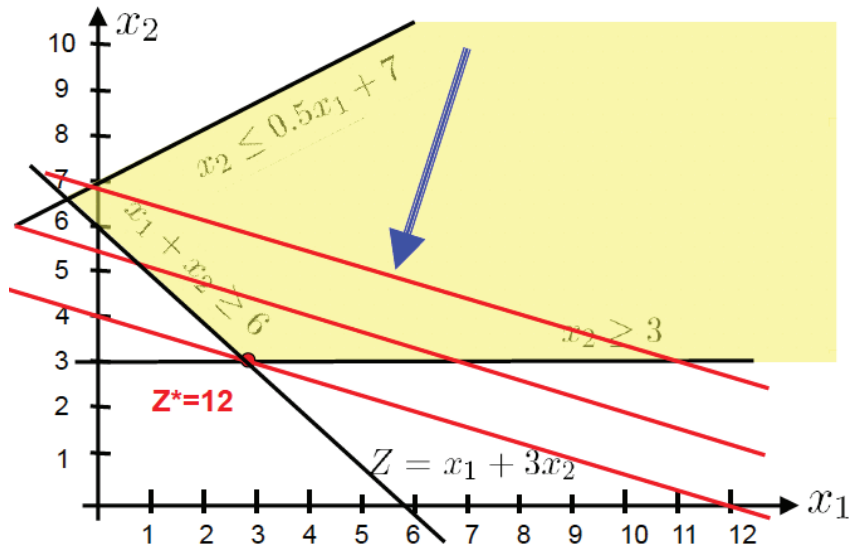
Feasible set is unbounded



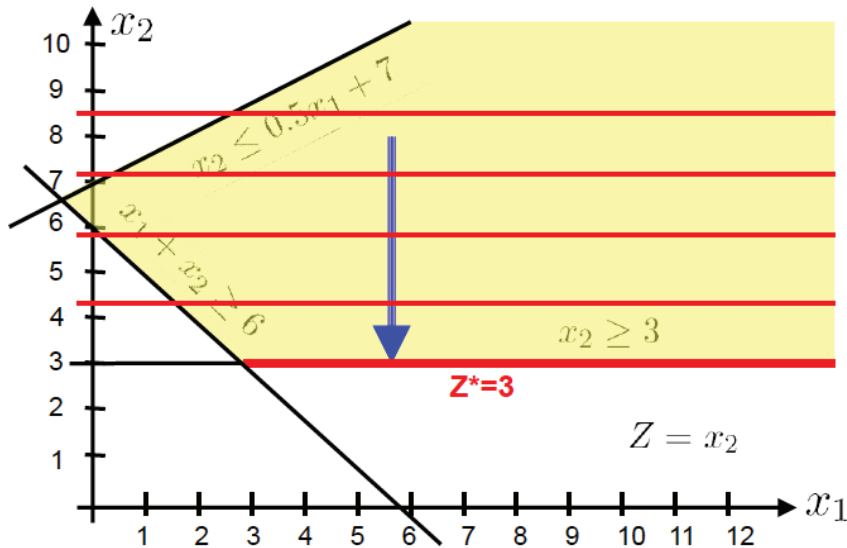
Objective function might be unbounded too



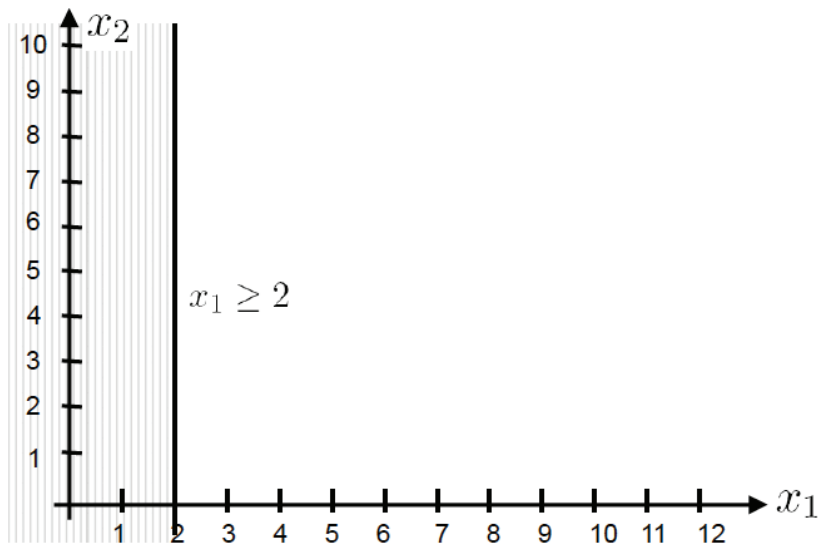
Objective function might be bounded



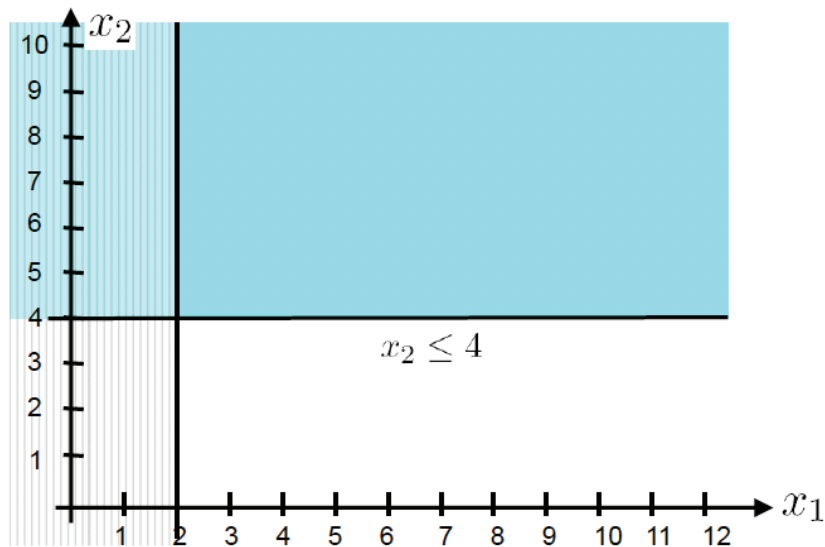
Optimum may be non-unique



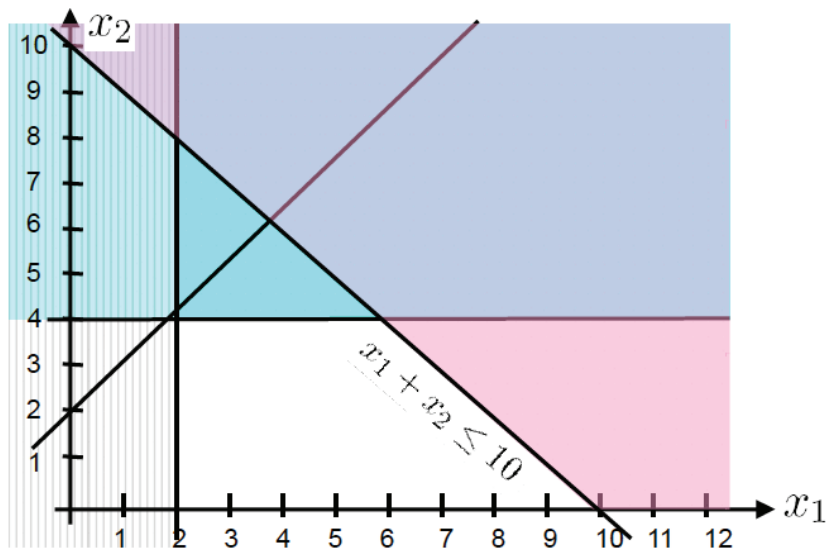
Feasible set might be empty



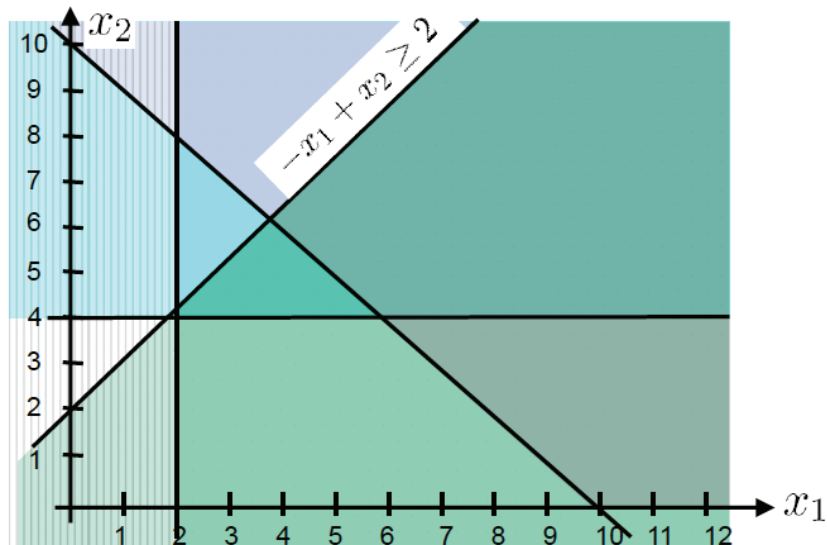
Feasible set might be empty



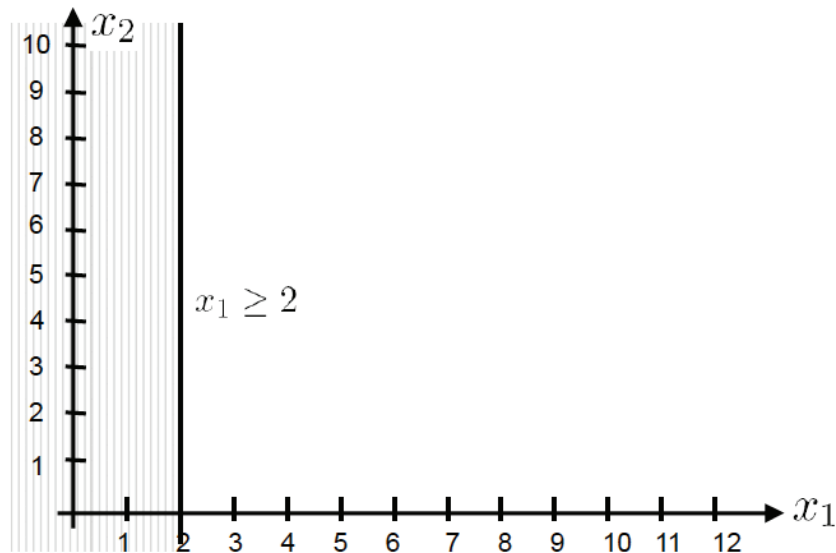
Feasible set might be empty



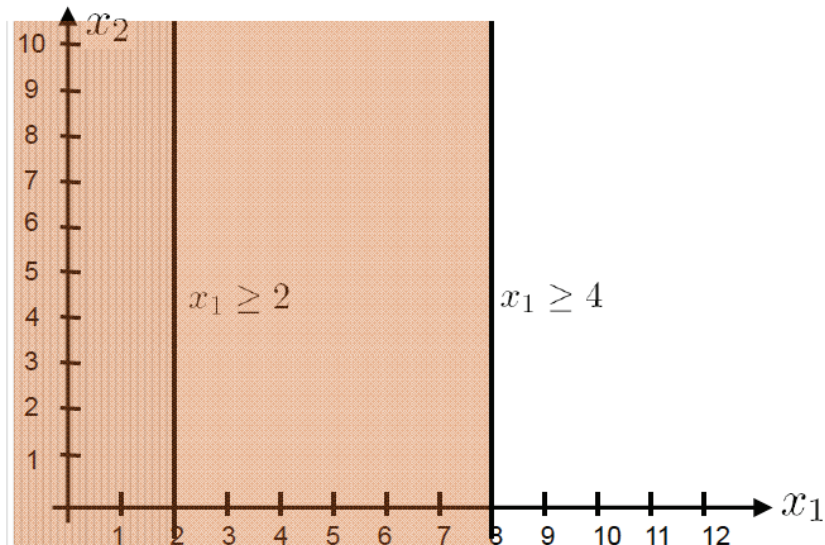
Feasible set might be empty



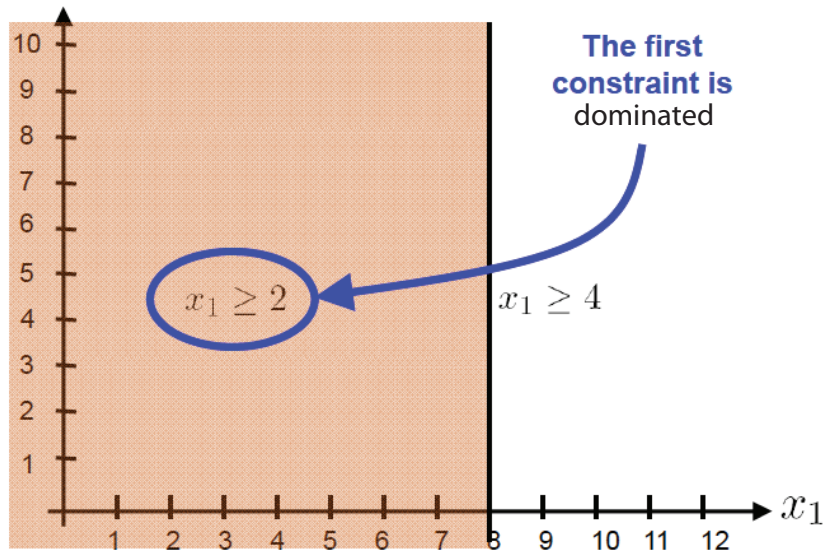
Constraint domination



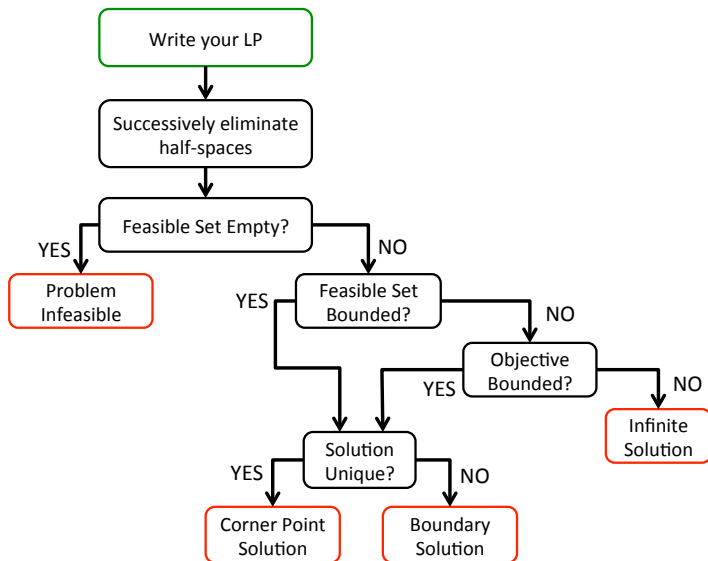
Constraint domination



Constraint domination

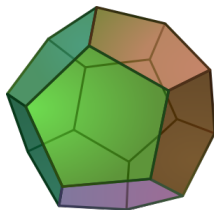


Graphical solution of LPs: A General Method



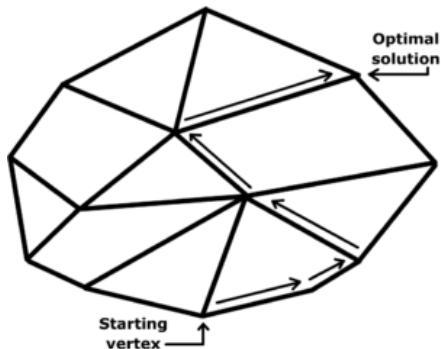
Insights from Graphical LP

- Linear constraints $Ax \leq b$ form feasible set (possibly empty)
- Feasible set is a (possibly unbounded) convex polytope
- Optimal solution exists along edges (corner point or line segment)



Danzig's Simplex Algorithm

- 1 Define feasible set
- 2 Start at vertex. Move along vertices until obj. fcn. stops decreasing



Example of Simplex Algorithm

Recall the LP problem:

$$\max \quad J = 140x_1 + 160x_2$$

$$\text{s. to} \quad 2x_1 + 4x_2 \leq 28$$

$$5x_1 + 5x_2 \leq 50$$

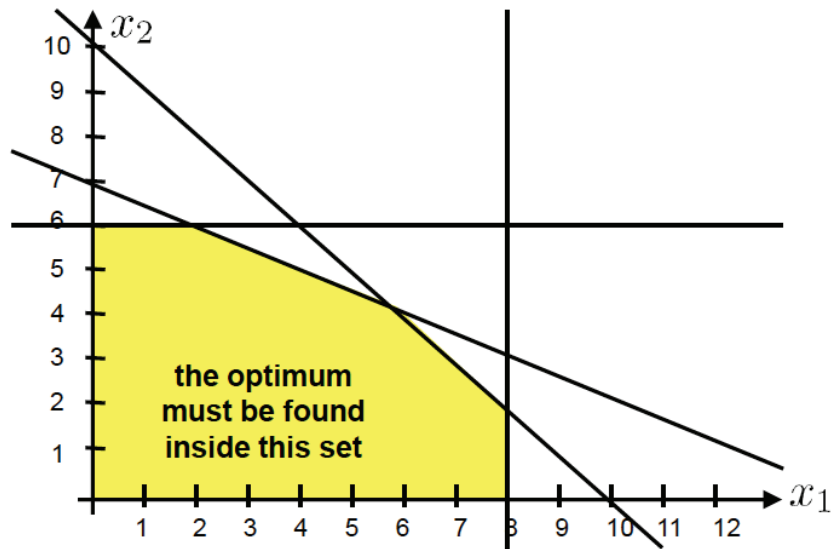
$$x_1 \leq 8$$

$$x_2 \leq 6$$

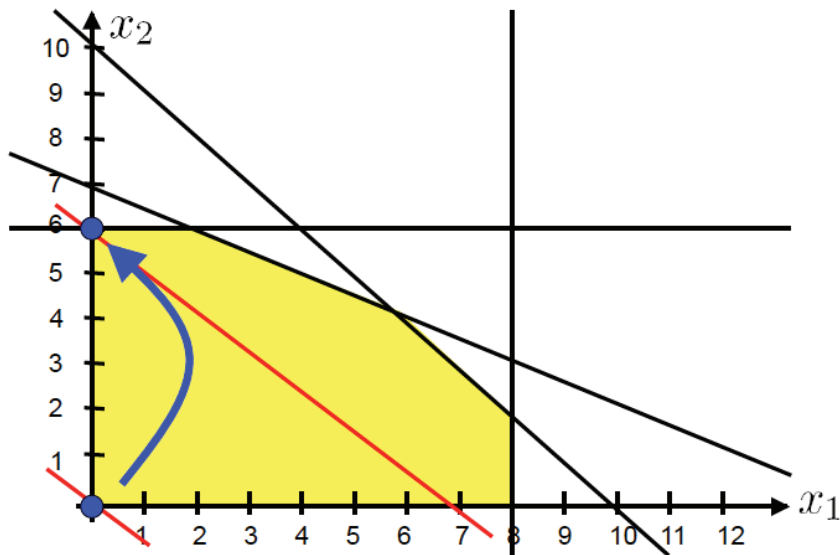
$$x_1 \geq 0$$

$$x_2 \geq 0$$

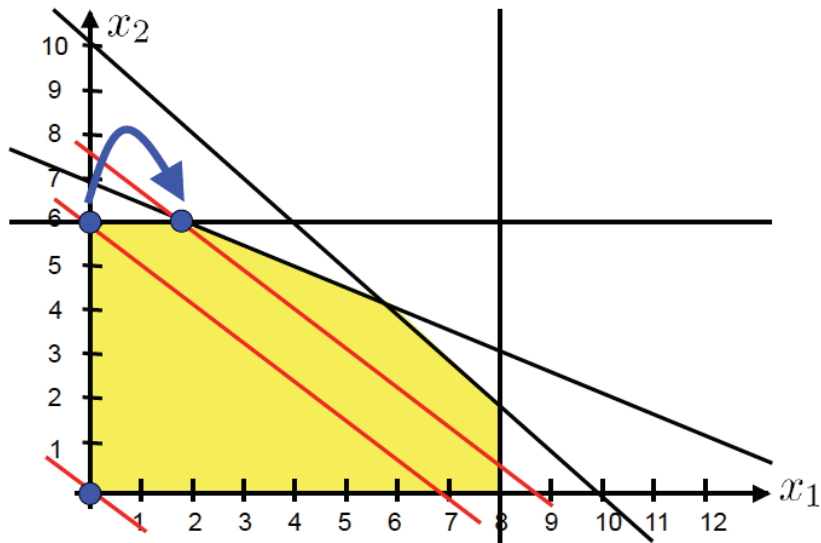
Feasible Set Final Result



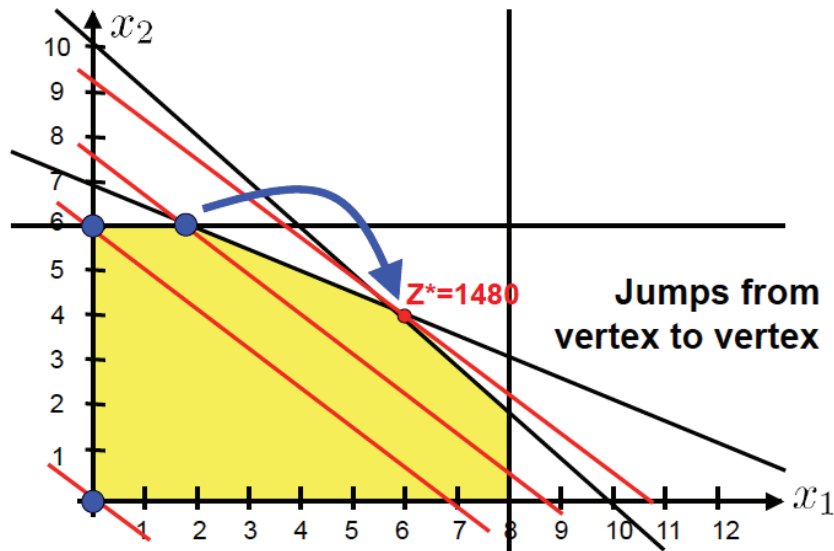
Start at a Vertex



Jump to adjacent vertex



Stop when objective stops decreasing



Revelle

- Chapter 3 - A Graphical Solution Procedure and Further Examples

Simplex Algorithm

- Revelle Chapter 4 - The Simplex Algorithm for Solving Linear Programs